

1. RATIONALITY, ALGEBRAICITY, AND HOLONOMICITY

We begin with an easy lemma.

Lemma 1.1. Let $f(x) = \sum a_n x^n \in \mathbf{Z}[[x]]$. Suppose that:

- (1) $f(x)$ is holomorphic in a disc of radius $R > 1$.
- (2) $a_n \in \mathbf{Z}$.

Then $f(x)$ is a polynomial.

Proof. We have the following estimates on a_n .

- (1) A Cauchy upper bound on a_n : By Cauchy's integral formula,

$$a_n = \frac{1}{2\pi i} \oint \frac{f(x)dx}{x^{n+1}}.$$

Take the contour to be a disc of radius η with $1 < \eta < R$ for some fixed η . The function $f(x)$ is uniformly bounded on that disc, say by a constant C . Hence

$$|a_n| \leq \frac{1}{2\pi} \oint_{|x|=\eta} \left| \frac{f(x)}{x^{n+1}} \right| = \frac{2\pi\eta}{2\pi} \frac{C}{\eta^{n+1}} = \frac{C}{\eta^n} = o(1).$$

- (2) A Liouville lower bound on a_n : If $f(x)$ is not a polynomial, then $a_n \neq 0$ for infinitely many integers n . For those integers n , we have the bound

$$|a_n| \geq 1.$$

We call this Liouville because in Liouville's estimate on $|\alpha - p/q|$ for a root α of a degree n polynomial, the key input is that $|f(p/q)q^n| \geq 1$ if it is non-zero.

These bounds are in conflict, hence $f(x)$ must be a polynomial. □

How can one improve on these theorems?

- (1) I would like to sketch improvements by Borel, Dwork, Polya, Betrandias,
- (2) Vesselin will talk about André's results — introduce our improvement, and explain the new ingredients missing from André.

1.1. Borel. In the trivial lemma above, let's drop the hypothesis that $f(x)$ is analytic and weaken it to being meromorphic instead.

Theorem 1.2 (Borel). Let $f(x) = \sum a_n x^n \in \mathbf{Z}[[x]]$. Suppose that:

- (1) $f(x)$ is meromorphic in a disc of radius $R > 1$.
- (2) $a_n \in \mathbf{Z}$.

Then $f(x)$ is a rational function.

Here's an incorrect proof. By compactness, $f(x)$ has only finitely many poles in a closed disc of radius η for $1 < \eta < R$. Let $q(x)$ be a polynomial which vanishes at all these points to the correct order — now apply the trivial lemma to $f(x)q(x)$ to deduce that $f(x)q(x) = p(x)$. The problem is that this only works if we can take $q(x) \in \mathbf{Z}[x]$, and this only works if the poles are algebraic numbers. But this doesn't have to be true a priori (but it does a posteriori).

The proof idea (all the proof ideas) are to combine an upper and lower bound. For rationality proofs, one strategy is to find a criteria strictly in terms of integers for when a

function $f(x)$ is a rational function. We know that coefficients of rational functions satisfy linear recursion relations (maybe with transcendental coefficients). It turns out that the key idea is to look at the determinants

$$\begin{vmatrix} a_k & a_{k+1} & \dots & a_{k+n} \\ a_{k+1} & a_{k+2} & \dots & a_{k+1+n} \\ & \ddots & & \\ a_{k+n} & a_{k+n+1} & \dots & a_{k+2n-1} \end{vmatrix}.$$

The vanishing of these determinants detects whether $f(x)$ is rational or not. For example, the coefficients a_n form an arithmetic progression if and only if the 3×3 determinants vanish; remember back to when you teach linear algebra and write down a random matrix to invert:

$$\begin{bmatrix} 1 & 2 & 3 \\ 2 & 3 & 4 \\ 3 & 4 & 5 \end{bmatrix}$$

and then find to your horror that it has determinant zero!

Let me explain the argument the case where there is a single pole at some transcendental point λ where $|\lambda| < 1$. (or $\lambda \leq 1$ as well. The case $|\lambda| > 1$ already follows by the trivial lemma). We have bounds

$$|a_n| = O(1/\lambda^n)$$

which could be big. But if $\sum b_n = f(x)(x - \lambda) = \sum x^n(\lambda a_n - \lambda a_{n-1})$, then

$$|b_n| = O(1/\eta^n).$$

is small. And now

$$\begin{aligned} \begin{vmatrix} a_k & a_{k+1} & \dots & a_{k+n} \\ a_{k+1} & a_{k+2} & \dots & a_{k+1+n} \\ & \ddots & & \\ a_{k+n} & a_{k+n+1} & \dots & a_{k+2n-1} \end{vmatrix} &= \begin{vmatrix} a_k & a_{k+1} & \dots & a_{k+n} \\ a_{k+1} - \lambda a_k & a_{k+2} - \lambda a_{k+1} & \dots & a_{k+1+n} - \lambda a_{k+n} \\ & \ddots & & \\ a_{k+n} - \lambda a_{k+n-1} & a_{k+n+1} - \lambda a_{k+n} & \dots & a_{k+2n-1} - \lambda a_{k+2n-2} \end{vmatrix} \\ &= \begin{vmatrix} a_k & a_{k+1} & \dots & a_{k+n} \\ b_k & b_{k+1} & \dots & b_{k+n} \\ & \ddots & & \\ b_{k+n-1} & b_{k+n} & \dots & b_{k+2n-2} \end{vmatrix}. \end{aligned}$$

If n is big enough with respect to η and λ , this is both very small and an integer, hence zero.

1.2. Dwork. The above argument work exclusively at the complex place, but there is a p -adic version as well, and versions which mix primes. Note however that any rational function in $\mathbf{Q}[[x]]$ — indeed any algebraic function, has coefficients in $\mathbf{Z}[1/S]$ for some integer S .

1.3. Polya–Bertrandias. Here we weaken the hypothesis that $f(x)$ extends to a meromorphic function for $R > 1$, but we still assume *some* convergence outside the disc. To give a completely concrete example, suppose that $f(x) \in \mathbf{Z}[[x]]$ extended to a holomorphic function on $\mathbf{C} \setminus [1, \infty)$. We can't use Cauchy on any disc of radius $\eta > 1$, so where do we get a good upper bound from?

Lemma 1.3. The only such functions have the form $P(x)/(1-x)^N$ for some integer N .

Proof. Consider

$$\oint \left(\frac{1}{x^2} - \frac{1}{x} \right)^n \frac{f(x)}{x} dx, \quad \oint \left(\frac{1}{x^2} - \frac{1}{x} \right)^n f(x) dx.$$

These integrals vanish for n large enough. The reason is that if we pick a contour close enough to the boundary (a fixed contour), we see that the function $f(x)$ is bounded, but the function $(1/x^2 - 1/x)$ is bounded above by $(1/4 + \varepsilon)$, and so this is very small. But it is also an integer, and hence it must be zero for n large enough.

After replacing $f(x)$ by $f(x)(1/x - 1/x^2)^N \in \mathbf{Z}[[x]](1/x)$ for large enough N , the integrals above vanish for all $n \in \mathbf{N}$.

We can write

$$\begin{aligned} \oint \frac{f(x)}{x^{2n+1}} &= \oint \left(\frac{1}{x^2} - \frac{1}{x} \right)^n \frac{f(x)}{x} dx + \sum_{i < 2n+1} c_{2n+1,i} \oint \frac{f(x)}{x^i}, \\ \oint \frac{f(x)}{x^{2n}} &= \oint \left(\frac{1}{x^2} - \frac{1}{x} \right)^n f(x) dx + \sum_{i < 2n} c_{2n,i} \oint \frac{f(x)}{x^i}, \end{aligned}$$

and hence by induction deduce that $f(x) \in \mathbf{Z}[1/x]$, and thus

$$f(x) = \frac{P(x)}{(1-x)^N}$$

for some polynomial $P(x)$. □

This proof works as long as $1/x^2 - 1/x$ has absolute value less than one near the boundary of the region, or equivalently that $x - x^2$ has absolute value less than one near the interval $[0, 1]$ after replacing x by $1/x$. This means that exactly the same argument works for $\mathbf{C} \setminus [\alpha, 0]$ where

$$\alpha > \frac{1}{\phi} = \frac{\sqrt{5} - 1}{2} = 0.618 \dots$$

because $|x - x^2| < 1$ on $[0, 1/\alpha]$ for $1/\alpha < \phi$. This is optimal, because

$$\frac{1}{1 - x - x^2} \in \mathbf{Z}[[x]]$$

But you can replace $x - x^2$ by

$$(x - x^2)(x - x^2 + 1)$$

and get a larger interval, and hence deduce rationality for even smaller α . So it's all about finding monic polynomials $f(x)$ which are less than 1 on the interval $[0, 1/\alpha]$, which relates (almost by definition) to the transfinite diameter of this interval $[0, 1/\alpha]$. Since $\mathbf{C} \setminus [\alpha, \infty)$

is simply connected, this can also be directly related to the conformal radius of the region. In particular, make the substitution

$$z \mapsto \frac{-4z}{(1-z)^2}$$

which maps $D(0, 1)$ to $\mathbf{C} \setminus [1, \infty)$, and realizes this region with conformal radius 4. If you replace $\mathbf{C} \setminus [1, \infty)$ by $\mathbf{C} \setminus [\alpha, \infty)$ then the conformal radius becomes 4α . So this works for $\alpha > 1/4$, and more generally:

Theorem 1.4 (Polya–Bertrandias). *Suppose that $f(x)$ extends as a meromorphic function to a region of conformal radius $R > 1$. Then $f(x)$ is rational.*

(Vesselin tells me that this is Polya’s theorem and Bertrandias’ contribution was the p -adic extension which is not included here.)

1.4. A small cheat. Note that in our example above we used that the polynomial $x^2 - x$ had coefficients in \mathbf{Z} . Having small conformal radius only guarantees the existence of a monic polynomial with arbitrary coefficients which is small. Yet we really used integrality in our arguments to deduce that a small integer had to be zero. If we replace it by $P(x)$, we deduce that

$$\oint P(x)^n f(x) x^i dx$$

is very small for $i = 0, \dots, \deg(P) - 1$ and for all large enough n . But it doesn’t have to be zero because it’s not an integer! However, we can once again study the Hankel determinants rather than the coefficients themselves, and deduce that $f(x)$ is rational. (In effect, the version of the argument above when $x - x^2$ happens to have integral coefficients is a little like the case of Borel’s theorem where the poles are at algebraic points.)

The example above is optimal in that

$$\sqrt{1 - 4x} \in \mathbf{Z}[[x]].$$

is not rational.

In all honesty, I’m doubt I will get much further than this point, so the rest is (even more) sloppy than what came before

1.5. Algebraicity. Looking at the example $\sqrt{1 - 4x}$, it has a stronger property than being defined on $\mathbf{C} \setminus [1/4, \infty)$. Namely, it has an analytic continuation to the universal cover of $\mathbf{C} \setminus \{0, 1/4\} = \mathbf{P}^1 - \{0, \alpha, \infty\}$ for $\alpha = 1/4$. Once more, let us consider the special case of multivalued functions on $\mathbf{P}^1 - \{0, \alpha, \infty\}$ with coefficients in $\mathbf{Z}[[x]]$. Algebraicity implies some $\mathbf{Q}[x]$ -linear relationship between the powers $f(x)^n \in \mathbf{Z}[[x]]$.

1.6. The strategy. Form a $\mathbf{Z}[x]$ linear combination of these powers:

$$F(x) = \sum a_i(x) f^i(x)$$

such that:

- (1) The polynomials $a_i(x)$ have moderate degree
- (2) The sum $F(x)$ vanishes to high order at $x = 0$.

(3) Make the coefficients of $a_i(x)$ not too big.

If we can do this, then with $F(x) = \sum a_n x^n$ we can play the game with bounding a_n for the smallest non-zero coefficient. It's non-zero assuming that $f(x)$ is not algebraic, but then obtain a Cauchy bound.

There is tension between these conditions. For the first two, bounding the degree and imposing vanishing is simply a system of linear equations. If we have enough equations (high enough degree) we get a solution. However, those solutions might be very large. So how do we control these solutions and hence the coefficients of $a_i(x)$? The key point is to have many more variables than equations so the solution space is big — for example, take twice as many variables as equations. We still get vanishing to high degree, but we have much more flexibility in looking for a small solution. Indeed, the classic Siegel's lemma is exactly about finding small solutions to such underdetermined linear equations.

1.7. André, Limits of this theorem.

Theorem 1.5. *If $f(x)$ is a multivalued function on $\mathbf{P}^1 \setminus \{0, \alpha, \infty\}$ which is holomorphic at (the original branch of) 0, and $\alpha > 1/16$, then $f(x)$ is algebraic.*

Actually, this is best possible. Take

$$F(1/2, 1/2, 1, 16x) = \sum \binom{2n}{n}^2 x^n \in \mathbf{Z}[[x]].$$

is a solution to an ODE in $\mathbf{P}^1 \setminus \{0, 1/16, \infty\}$. But I claim it is not an algebraic function. We saw that

$$F(1/2, 1/2, 1, 16\lambda/16) = F(1/2, 1/2, 1, \lambda) = (1 + 2q + 2q^4 + 2q^9 + \dots) = \theta^2(\tau)$$

Now $\theta^2(\tau)$ has weight one and is not an algebraic function of λ which has weight zero.

In this particular example, the monodromy around 0 (of the other solution!) is unipotent. The solutions are $\{1, \tau\}$, and around a loop we should get $\tau + 2$, since that doesn't change $q = e^{\pi i \tau}$. The monodromy is actually $\Gamma(2)$ since it is the Picard–Fuchs equation.