PCMI GSS Asymptotic Enumeration: Problem Set 1

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The \star 's roughly indicate difficulty. Hints are given as footnotes. Please find a group to work with and don't be afraid to ask questions!

- 1. \star Recall this theorem from the lecture: if G is an n-vertex d-regular bipartite graph, then $\log_2 i(G) \leq (1 + O(1/d)) n/2$. Give an example which shows that the dependence on d is best possible. ($\star\star$ If we've prove this theorem already, try to follow the proof on your example and notice how the inequalities become sharp.)
- 2. $\star \star \star$ You might skim through this exercise if you are already familiar with Shannon entropy. In what follows $\mathbf{X}, \mathbf{Y}, \ldots$ are finitely-supported random variables. Recall that the (binary) entropy of \mathbf{X} is

$$H(\mathbf{X}) = \sum_{x} \mathbb{P}(\mathbf{X} = x) \log \frac{1}{\mathbb{P}(\mathbf{X} = x)},$$

where the log is taken base 2. Entropy is a measure of the 'information' stored in a random variable, here measured in bits; another interpretation is that entropy is the average 'surprise' upon learning the realization of a random variable. Verify the following starting with only the above definition of entropy.

- (a) (Uniform maximizes entropy) Suppose **X** is supported on a finite set S. Show that $H(\mathbf{X}) \leq \log |S|$, with equality if and only if **X** is uniform on S. (This is why entropy is helpful for counting problems: if we want to count |S|, we may equivalently calculate the entropy of the uniform random variable on S.)¹
- (b) (Chain rule) Recall the definition of conditional entropy:

$$H(\mathbf{X}|\mathbf{Y}) = \sum_{y} \mathbb{P}(\mathbf{Y} = y) \sum_{x} \mathbb{P}(\mathbf{X} = x | \mathbf{Y} = y) \log \frac{1}{P(\mathbf{X} = x | \mathbf{Y} = y)}.$$

Show that

$$H(\mathbf{X}, \mathbf{Y}) - H(\mathbf{X}) = H(\mathbf{Y}|\mathbf{X}).$$

¹There is an elementary proof using the inequality $\log(x) \leq x - 1$. Alternatively, one can use Jensen's inequality.

- (c) (Additivity for independent variables) We say that \mathbf{X} and \mathbf{Y} are independent if $\mathbb{P}(\mathbf{X} = x, \mathbf{Y} = y) = \mathbb{P}(\mathbf{X} = x)\mathbb{P}(\mathbf{Y} = y)$ for all x and y. Show that if \mathbf{X} and \mathbf{Y} are independent, then $H(\mathbf{Y}|\mathbf{X}) = H(\mathbf{Y})$ and $H(\mathbf{X}, \mathbf{Y}) = H(\mathbf{X}) + H(\mathbf{Y})$.
- (d) (Dropping conditioning) Show that $H(\mathbf{Y}|\mathbf{X}) \leq H(\mathbf{Y})$, and characterize equality.²
- (e) (Subadditivity) Show that $H(\mathbf{X}, \mathbf{Y}) \leq H(\mathbf{X}) + H(\mathbf{Y})$. Furthermore $H(\mathbf{X}) \leq \sum_{i} H(\mathbf{X}_{i})$.
- (f) (Dropping conditioning, part 2) Show that $H(Y|X, Z) \le H(Y|X)$.
- (g) (Inserting conditioning) Show that $H(\mathbf{X}|\mathbf{Z}) \leq H(\mathbf{X}|\mathbf{Y}) + H(\mathbf{Y}|\mathbf{Z})$.
- (h) (Refinement) If Y determines X, then H(X|Y) = 0 and $H(Z|Y) \le H(Z|X)$.
- (i) (Shearer's Inequality) Let $\mathbf{X} = (\mathbf{X}_1, \dots, \mathbf{X}_k)$ be a random vector, and let $\alpha : 2^{[k]} \to \mathbb{R}_{\geq 0}$ be such that $\sum_{A \ni i} \alpha(A) \geq 1$ for all $i \in [k]$. Show that

$$H(\mathbf{X}) \le \sum_{A \subseteq [k]} \alpha(A) H(\mathbf{X}_A),$$

where $\mathbf{X}_A = (\mathbf{X}_i)_{i \in A}$.

3. ** Let $0 \le \alpha \le 1/2$. Show that the number of subsets of [n] of size at most αn is at most $2^{h(\alpha)n}$, where $h(\alpha) = \alpha \log_2\left(\frac{1}{\alpha}\right) + (1-\alpha) \log_2\left(\frac{1}{1-\alpha}\right)$. Derive that $\binom{n}{\le k} \le \left(\frac{en}{k}\right)^k$.

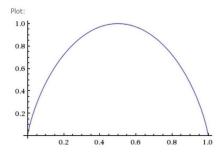


Figure 1: a plot of $h(\alpha)$

- 4. $\star\star$ Find a bijection between the following two collections.
 - (proper) 3-colorings of Q_d with the color of a prefixed vertex v_0 is fixed;
 - the graph homomorphisms from Q_d to \mathbb{Z} (in which two vertices a, b are adjacent iff |a b| = 1) with $f(v_0) = 0$.

²Use $\log(x) \le x - 1$ or Jensen's Inequality *judiciously*, in addition to Bayes' Rule.

³For each $A \subseteq [k]$, we have $H(\mathbf{X}_A) = \sum_{i \in A} H(\mathbf{X}_i | \mathbf{X}_{A \cap [i-1]}) \ge \sum_{i \in A} H(X_i | X_{[i-1]})$.

⁴Image from D. Galvin, Three tutorial lectures on entropy and counting, arXiv 1406.7872

⁵Where have you seen $h(\alpha)$ before? Express the counting problem as an entropy problem and use subadditivity.