

Lecture 1

Motivating examples

- $\mathcal{I}(G) := \{ \text{indep sets in } G \}$, $i(G) := |\mathcal{I}(G)|$
vt set in which no pairs are adj.

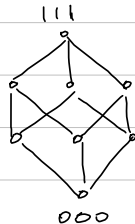
E.g. 1. G : d -reg, n vts, bip. (default: d, n large)
_____ $n/2$
_____ $n/2$ $i(G) \geq 2^{n/2}$

Prop. $i(G) \leq 2^{(1+o_d(1))n/2}$ ($= 2^{n/2} \times 2^{o(n/2)}$)

Pf. entropy, graph container (main focus)

E.g. 2. Q_d : d -dimensional hypercube

e.g. Q_3



$$V(Q_d) = \{0, 1\}^d$$

$v \sim w$ iff differ on exactly
1 word

Q_d : d -reg, bip on $\mathbb{E} \cup \emptyset$ (unique) _____ $\mathbb{E}$

$$|V(Q_d)| = 2^d =: N$$

_____ $\emptyset$

$i(G)?$

Thm. (Korshunov - Sapozhenko '83)

$$i(Q_d) \sim 2 \cdot \sqrt{e} \cdot 2^{M/2}$$

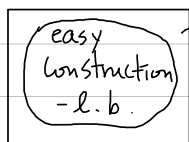
as $d \rightarrow \infty$

$$f \sim g \Leftrightarrow \frac{f}{g} \rightarrow 1$$

i.e., $f = (1 + o(1))g$

No closed formula for $i(Q_d)$.

Pf goes by



negli.

→ typical structure

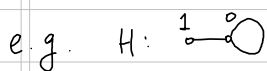
(e.g. sampling arg.)

$I(Q_d)$

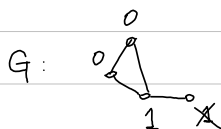
IS $\xrightarrow{\text{gen.}}$ graph hom's.

Def. $f: V(G) \rightarrow V(H)$ is a graph homomorphism if

$$\{u, v\} \in E(G) \Rightarrow \{f(u), f(v)\} \in E(H)$$

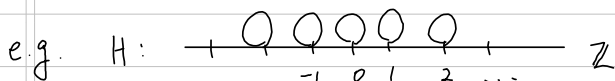
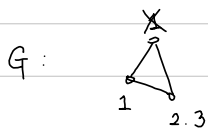


IS

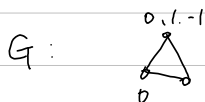


e.g. proper coloring

H : K_3



1-Lipschitz function



Entropy basics (Pfs: EX)

X, Y : discrete rand var's

$$p(x) := P(X=x) \quad , \quad \log = \log_2$$

Entropy of X : $H(X) = \sum_x p(x) \log \frac{1}{p(x)}$ $\leadsto$ amount of
(randomness (surprise, info) in X)
($0 \cdot \log \frac{1}{0} = 0$)

Properties:

$$\textcircled{1} H(X) \leq \log |\text{Range}(X)|, \quad = \text{iff } X \text{ uniform } \quad \text{"most random"}$$

↑ connection

$$(\because) H(X) = \sum_x p(x) \cdot \log \frac{1}{p(x)} \leq \log \left(\sum_x p(x) \cdot \frac{1}{p(x)} \right) \quad \square$$

↑
Jensen

$$\textcircled{2} H(X_1, \dots, X_k) \leq H(X_1) + \dots + H(X_k)$$

$$\text{Want(?)} H(X, Y) = H(X) + \underbrace{H(Y|X)}_{\text{conditional entropy}}$$

Def. T : event,

$$H(X|T) = \sum_x p(x|T) \log \frac{1}{p(x|T)} \quad p(x|T) = P(X=x|T)$$

Y : r.v.

$$\begin{aligned} H(X|Y) &= \mathbb{E}[H(X|Y=y)] \\ &= \sum_y p(y) \cdot H(X|Y=y) \end{aligned}$$

$$\textcircled{3} \quad H(X_1, X_2, \dots, X_k) = H(X_1) + H(X_2|X_1) + H(X_3|X_1, X_2) + \dots$$

$$\textcircled{4} \quad \text{If } Y \text{ determines } Z, \text{ then} \\ H(X|Y) \leq H(X|Z)$$

Shearer's Lemma gen of $\textcircled{2}$

$$\text{e.g. } X = (X_1, X_2, X_3, X_4)$$

$$H(X) \leq \frac{2}{3} H(X_1, X_2, X_3) + \frac{1}{3} [H(X_1, X_4) + H(X_2, X_4) + H(X_3, X_4)]$$

Lem $X = (X_1, \dots, X_k)$ random vector, $\alpha: 2^{[k]} \rightarrow \mathbb{R}^{>0}$

If $\sum_{A \ni i} \alpha(A) \geq 1 \quad \forall i \in [k]$, then

$$H(X) \leq \sum_{A \subseteq [k]} \alpha(A) \cdot H(\underbrace{X_A}_{= (X_i : i \in A)})$$

Prop. G : d -reg, n vts, bip.

$$i(G) \leq 2^{(1+O(1/d))n/2}$$

starting pt: $f \in \mathcal{I}(G)$ unif rand $\Rightarrow H(f) = \log i(G)$

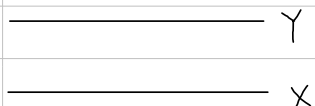
$$\text{ETS. } H(f) < (1+O(1/d))n/2$$

$$\boxed{\text{ETS. } H(f) < (1 + O(\frac{1}{d})) \frac{n}{2}}$$

view $f: V(G) \rightarrow \{0, 1\}$

notation: $f_x = f(x)$, for $A \subseteq V$, $f_A = (f_x : x \in A)$

$f(A) = \{f_x : x \in A\}$ "palette"



$N_x = N(x)$

$\hookrightarrow \{0\} \text{ or } \{1\} \text{ or } \{0, 1\}$

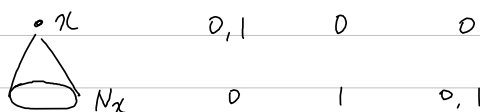
$$H(f) = H(f_Y) + H(f_X | f_Y) \quad (*)$$

$$H(f_X | f_Y) \leq \sum_{x \in X} H(f_x | f_Y) \leq \sum_{x \in X} H(f_x | f(N_x))$$

$$H(f_Y) \stackrel{\text{Sh}}{\leq} \frac{1}{d} \sum_{x \in X} H(f_{N_x}) = \frac{1}{d} \sum_{x \in X} [H(f(N_x)) + H(f_{N_x} | f(N_x))]$$

$$\rightarrow (*) \leq \sum_{x \in X} \underbrace{\frac{1}{d} H(f(N_x))}_{O(1)} + \underbrace{[H(f_x | f(N_x)) + \frac{1}{d} H(f_{N_x} | f(N_x))]}_{(*)} \quad \text{Claim. } \leq 1 \quad (\text{then done})$$

picture:



rigorously:

$$(*) = \sum_c P(f(N_x) = c) \underbrace{[H(f_x | f(N_x) = c) + \frac{1}{d} H(f_{N_x} | f(N_x) = c)]}_{\leq 1} \leq 1 \quad \square$$