## SHEET 4

## 1 Relation between maximum and average error probability

Show that for any M-code (E,D) for a channel  $\{W_x\}$ , there exists an  $\lfloor M/2 \rfloor$ -code (E',D') for the same channel with  $p_{\text{err,max}}(E',D') \leq 2p_{\text{err}}(E,D)$ .

## 2 Channel capacities

- 1. For a classical-quantum channel  $\{W_x\}_{x\in\mathcal{X}}$ , we defined the function  $f(n)=\sup_{P_{X^n}}I(X^n:B^n)_{\rho}$  where  $\rho_{X^nB^n}=\sum_{x^n\in\mathcal{X}^n}P_{X^n}(x^n)|x^n\rangle\langle x^n|_{X^n}\otimes W_{x^n}$  and  $W_{x^n}=W_{x_1}\otimes W_{x_2}\cdots W_{x_n}$ . Show that the function f is superadditive.
- 2. Let  $p \in [0,1]$  and consider the channel W with input alphabet  $\mathcal{X} = \{0,1\}$  defined by  $W_0 = (1-p)|0\rangle\langle 0| + p|1\rangle\langle 1|$  and  $W_1 = p|0\rangle\langle 0| + (1-p)|1\rangle\langle 1|$  (this is called the binary symmetric channel). Compute the capacity C(W).
- 3. Let  $e \in [0,1]$  and consider the channel W with input alphabet  $\mathcal{X} = \{0,1\}$  defined by  $W_0 = (1-e)|0\rangle\langle 0| + e|E\rangle\langle E|$  and  $W_1 = (1-e)|1\rangle\langle 1| + e|E\rangle\langle E|$ , where  $|E\rangle$  is orthogonal to both  $|0\rangle$  and  $|1\rangle$  (this is called the binary erasure channel). Compute the capacity C(W).
- 4. Prove that the von Neumann entropy is strictly concave, i.e., for any distribution  $\{p(x)\}$  with support S we have  $H(\sum_{x\in S}p(x)W_x)\geq \sum_{x\in S}p(x)H(W_x)$  with equality if and only if  $W_x=W$  for all  $x\in S$ . (You may want to use a question from Sheet 3 for this). Conclude that a classical-quantum channel  $\{W_x\}_{x\in \mathcal{X}}$  satisfies C(W)=0 if and only if there exists W such that  $W_x=W$  for all  $x\in \mathcal{X}$ .

## 3 Properties of the Holevo information

Recall that the Holevo information  $\chi(\{p(x), \sigma_A^x\})$  of an ensemble  $\{p(x), \sigma_A^x\}$  is defined as  $I(X:A)_\rho$  where  $\rho_{XA} = \sum_x p(x)|x\rangle\langle x|_X \otimes \sigma_A^x$ .

- 1. Show that  $\chi(\{p(x), \sigma_A^x\}) = H(\sum_x p(x)\sigma_A^x) \sum_x p(x)H(\sigma_A^x)$
- 2. Show that  $\chi$  is concave in  $\{p(x)\}$  (for fixed  $\{\sigma_A^x\}$ ) and convex in  $\{\sigma_A^x\}$  (for fixed  $\{p(x)\}$ ).
- 3. The Holevo information of a channel  $\mathcal{W}: \mathrm{L}(A) \to \mathrm{L}(B)$  is defined as the supremum over all ensembles  $\{p(x), \sigma_A^x\}$  of the Holevo information  $\chi(\{p(x), \mathcal{W}(\sigma_A^x)\})$ . Show that we may restrict the optimization to ensembles where  $\sigma_A^x$  are all pure states.