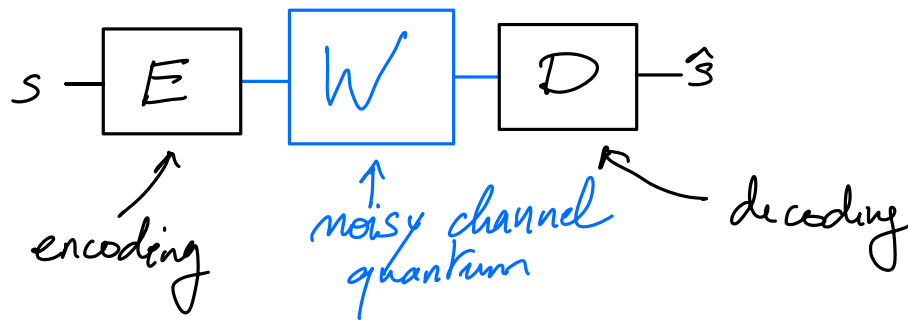


QUANTUM INFORMATION THEORY.



WANT : * $s \simeq \hat{s}$ with high probability
* maximize bits or qubits in s .

Objective : Understand fundamental limits

Plan : ① States, channels.

② State discrimination, Stein's lemma.

③ Data processing for the quantum relative entropy
Also called strong subadditivity of von Neumann entropy

④ Classical communication over quantum channels.
Shannon channel coding will be a special case.

⑤ Quantum communication over quantum channels.

Finite dimensional Hilbert space \mathcal{H} .

$u, v \in \mathcal{H}$ $\langle u, v \rangle$ inner product

$$\lambda \in \mathbb{C} \quad \langle u, \lambda v \rangle = \lambda \langle u, v \rangle$$

$$\langle \lambda u, v \rangle = \bar{\lambda} \langle u, v \rangle$$

\uparrow complex conjugate.

Linear operators $\mathcal{H} \rightarrow \mathcal{H}' : L(\mathcal{H}, \mathcal{H}')$

$$L(\mathcal{H}, \mathcal{H}) =: L(\mathcal{H})$$

For an operator $S \in L(\mathcal{H}, \mathcal{H}')$, the adjoint S^* is defined by

\uparrow
 $L(\mathcal{H}', \mathcal{H})$

$$\langle u', Su \rangle = \langle S^* u', u \rangle \text{ for all } u \in \mathcal{H}, u' \in \mathcal{H}'.$$

Important classes of operators $S \in L(\mathcal{H})$:

• S is unitary if $SS^* = S^*S = I$

• S is Hermitian if $S^* = S$. \uparrow identity.

$S \geq 0$
• S is positive, we write $S \in \text{Pos}(\mathcal{H})$ if S is Hermitian and $\langle u, Su \rangle \geq 0$ for all $u \in \mathcal{H}$.

• S is an orthogonal projection if $S^2 = S = S^*$ such as S is positive.

Bra-ket notation:

We identify $u \in \mathcal{H}$ with $|u\rangle \in L(\mathbb{C}, \mathcal{H})$

defined by $|u\rangle : \mathbb{C} \rightarrow \mathcal{H}$ \uparrow ket

$$\lambda \mapsto \lambda \cdot u.$$

The adjoint $|u\rangle^* \in L(\mathcal{H}, \mathbb{C})$ is denoted $\langle u|$ \leftarrow bra

$$\langle u | : \mathcal{H} \rightarrow \mathbb{C}$$

$$v \mapsto \langle u, v \rangle$$

We have $\cdot \langle u, v \rangle \in L(\mathbb{C}, \mathbb{C})$ identified with \mathbb{C} .
 $\langle u, v \rangle$.

\Rightarrow will denote inner product by $\langle u | v \rangle$.

$$\cdot |v\rangle \langle u| \in L(\mathcal{H})$$

Ex: e_i is a basis of \mathcal{H} then
 \uparrow orthonormal
then $I = \sum_i |e_i\rangle \langle e_i|$.

Rk: We will often use shorthand $|i\rangle$ for $|e_i\rangle$
 $|x\rangle$ for $|e_x\rangle$
 \vdots

Spectral decomposition

For any Hermitian $S \in L(\mathcal{H})$, there exists an orthonormal basis of \mathcal{H} $\{|e_i\rangle\}$ s.t.

$$S = \sum_{i=1}^{\dim(\mathcal{H})} \lambda_i |e_i\rangle \langle e_i|$$

with $\lambda_i \in \mathbb{R}$.

In other words S written in ONB $\{|e_i\rangle\}$ is diagonal

$$S = \begin{pmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_{\dim(\mathcal{H})} \end{pmatrix}$$

• S is positive iff $\lambda_i \geq 0 \quad \forall i$.

• For $f: \mathbb{R} \rightarrow \mathbb{C}$.

$$f(S) = \sum_i f(\lambda_i) |e_i\rangle\langle e_i|$$

Tensor products:

Multiple systems $A, B, C, \dots \quad X, Y, \dots$

Hilbert space $\begin{matrix} \downarrow & \downarrow & \downarrow \\ H_A & H_B & H_C \end{matrix}$

Hilbert space for joint system: $H_A \otimes H_B$ bilinear.

• vector space spanned by $u \otimes v$ for $u \in H_A, v \in H_B$

• inner product: $\langle u' \otimes v' | u \otimes v \rangle = \langle u' | u \rangle \cdot \langle v' | v \rangle$
and linear extension.

For $S \in L(H_A, H'_A), T \in L(H_B, H'_B)$ define $S \otimes T$:

$$(S \otimes T)(u \otimes v) = (Su) \otimes (Tv), \text{ and linear extension}$$

$\begin{matrix} \uparrow & \uparrow \\ H_A & H_B \end{matrix}$

$$= \text{span} \{ S \otimes T \}$$

We identify: $L(\mathcal{H}_A, \mathcal{H}'_A) \otimes L(\mathcal{H}_B, \mathcal{H}'_B)$ and $L(\mathcal{H}_A \otimes \mathcal{H}_B, \mathcal{H}'_A \otimes \mathcal{H}'_B)$

In particular $|u\rangle \otimes |v\rangle = |u \otimes v\rangle$

fixed on basis.

$$\text{Tr} S = \sum_i \langle e_i, S e_i \rangle$$

Def: A **density operator** ρ on \mathcal{H} is a normalized positive operator on \mathcal{H} , i.e., $\rho \in \text{Pos}(\mathcal{H})$ and $\text{Tr}(\rho) = 1$.

- The set of density operators is denoted $\mathcal{S}(\mathcal{H})$.
- ρ is said to be pure if $\text{rank}(\rho) = 1$
 $\rho \hat{=} |\psi\rangle\langle\psi| \quad \psi \in \mathcal{H}$.
- $\rho = \frac{1}{\dim \mathcal{H}} \cdot I$ "maximally" mixed.

Density operator formalism:

• If a system is represented by a vector $|\psi\rangle$ (e.g. $\alpha|0\rangle + \beta|1\rangle$) then the density operator ρ representing this system is given by $\rho = |\psi\rangle\langle\psi|$.

e.g. $\rho = \begin{pmatrix} |\alpha|^2 & \alpha\bar{\beta} \\ \bar{\alpha}\beta & |\beta|^2 \end{pmatrix}$ in the basis $\{|0\rangle, |1\rangle\}$.

• **Composition**: State of a composite system is given by density operators on $\mathcal{H}_A \otimes \mathcal{H}_B$ if individual state spaces are \mathcal{H}_A and \mathcal{H}_B .

Prepare ρ_A on system A and independently ρ_B on B: $\rho_A \otimes \rho_B$.

Notation: $\mathcal{H}_A \rightarrow A$.

(for short we call the Hilbert space A)

- Evolution: Isolated evolution of a subsystem A corresponds to a unitary on A. For a state ρ_{AB} on composite system $A \otimes B$ with evolution on A given by U_A and the B system unchanged:

$$\rho'_{AB} = (U_A \otimes I_B) \rho_{AB} (U_A^\dagger \otimes I_B).$$

- **Measurement**: A measurement on subsystem A is defined by operators $\{M_x\}_{x \in X}$ for some set X, $M_x \in \mathcal{L}(A)$ satisfying $\sum_{x \in X} M_x^\dagger M_x = I$

Prob of outcome x : $p(x) = \text{Tr}(M_x \otimes I_B \rho_{AB} M_x^\dagger \otimes I_B)$

$$\text{Check: } \sum_x p(x) = \sum_x \text{Tr}(M_x^\dagger M_x \otimes I_B \rho_{AB}) = \text{Tr}(\rho_{AB}) = 1.$$

\uparrow
 $\text{Tr}(ST) = \text{Tr}(TS).$

Post-measurement state conditioned on x :

$$\rho'_{AB, x} = \frac{(M_x \otimes I_B) \rho_{AB} (M_x^\dagger \otimes I_B)}{p(x)}.$$

• Special case: projective measurement

You might be used to special case

$M_x = P_x$ with P_x projector (ie $P_x^* = P_x = P_x^2$)
coming from the spectral decomposition of observable O

$$O = \sum_x \lambda_x \cdot P_x$$

A general measurement can model eg, a unitary followed by a projective measurement: U_A followed by $\{P_x\}_{x \in X}$

$$p(x) = \text{Tr} \left(\underbrace{P_x U_A}_{M_x} \otimes I_B \right) \rho_{AB} \left(\underbrace{U_A P_x}_{M_x^*} \otimes I_B \right)$$

• Special case: POVM measurement.

Often, we are not interested in post-measurement state but only in the probability distribution $p(x)$.

$$\begin{aligned} p(x) &= \text{Tr} \left(P_x \otimes I_B \right) \rho_{AB} \left(M_x^* \otimes I_B \right) \\ &= \text{Tr} \left(M_x^* M_x \otimes I_B \right) \rho_{AB} \end{aligned}$$

We let $E_x = M_x^* M_x$.

↗ only need to know E_x to determine $p(x)$ and not M_x .

Def: A positive operator valued measure (POVM) on A is a family $\{E_x\}_{x \in X}$ of positive operators on A such that $\sum_{x \in X} E_x = I_A$

Probability of outcome $x = \text{Tr}(E_x \rho)$.

Quantum channels

- General way of describing evolution of state of a system
- The Hilbert space can change: $A \rightarrow B$. (forget a system, add a particle...)
- \mathcal{E} should map $S(A)$ to $S(B)$

Def: A quantum channel \mathcal{E} is a **linear** map from $L(A)$ to $L(B)$ satisfying:

↑
maps convex combinations
 $\mathcal{E}(p\rho_0 + (1-p)\rho_1) = p\mathcal{E}(\rho_0) + (1-p)\mathcal{E}(\rho_1)$

- Completely positive:

For any Hilbert space R

$$\rho \in \text{Pos}(A \otimes R) \implies (\mathcal{E} \otimes I_R)(\rho) \in \text{Pos}(B \otimes R)$$

↑
Identity on $L(R)$: "superoperator"

- Trace-preserving:

$$\text{For } T \in L(A), \text{Tr}(\mathcal{E}(T)) = \text{Tr}(T)$$

Remark: Why complete positivity and not just positivity?

A positive map \mathcal{E} : For any $\rho \in \text{Pos}(A)$, $\mathcal{E}(\rho) \in \text{Pos}(B)$.

It turns out that positivity of a map is not stable under tensor product i.e.

$\exists \mathcal{E}, \mathcal{F}$ positive maps $\mathcal{E}: L(A) \rightarrow L(B)$, $\mathcal{F}: L(\bar{A}) \rightarrow L(\bar{B})$
such that $\mathcal{E} \otimes \mathcal{F}$ defined by:

$$(\mathcal{E} \otimes \mathcal{F})(S \otimes T) = \mathcal{E}(S) \otimes \mathcal{F}(T)$$

is **not** positive

Clearly if $\rho \in \text{Pos}(A)$, $\sigma \in \text{Pos}(\bar{A})$ then

$$(\mathcal{E} \otimes \mathcal{F})(\rho \otimes \sigma) \in \text{Pos}(B \otimes \bar{B})$$

$$\text{So } (\mathcal{E} \otimes \mathcal{F})(\text{Sep}(A; \bar{A})) \subseteq S(B \otimes \bar{B})$$

where $\text{Sep}(A; \bar{A}) = \text{conv} \left\{ \rho \otimes \sigma : \begin{array}{l} \rho \in S(A) \\ \sigma \in S(\bar{B}) \end{array} \right\}$
 \uparrow separable states.

$(\mathcal{E} \otimes \mathcal{F})(\rho)$ can fail to be positive for $\rho \in S(A \otimes \bar{A}) \setminus \text{Sep}(A; \bar{A})$

Typical example showing positivity is not stable:

\mathcal{E} : Transpose map

\mathcal{F} : identity map.

Complete positivity is stable under tensor product.

Ex: • $\mathcal{E}: L(A) \rightarrow L(A)$ U unitary in A

$$\mathcal{E}(T) = U T U^* \quad \text{for any } T \in L(A).$$

* Completely positive:

$$(\mathcal{E} \otimes I_R)(\rho) = (U \otimes I_R) \rho (U^* \otimes I_R) \geq 0$$

$$[\langle v, (U \otimes I_R) \rho (U^* \otimes I_R) v \rangle$$

$$= \langle (U^* \otimes I_R) v, \rho (U^* \otimes I_R) v \rangle \geq 0]$$

More generally, a map

$$\mathcal{E}(T) = S T S^* \quad \text{for all } T$$

is completely positive.

* Trace preserving: $\text{Tr}(U T U^*) = \text{Tr}(U^* U T) = \text{Tr}(T)$.

• Partial trace map.

$\rho_{AB} \in S(A \otimes B)$ state of a composite system.

What is the state of system A on its own?

→ Should be a valid quantum channel, corresponds to "forgetting" B .

$$\text{Tr}_B: L(A \otimes B) \rightarrow L(A)$$

$$T \mapsto \sum_b (I_A \otimes \langle b|) T (I_A \otimes |b\rangle)$$

where $\{|b\rangle\}_b$ forms a basis of B .

Note that $\text{Tr}_B = \underbrace{I_A}_{\text{identity: } L(A) \rightarrow L(A)} \otimes \underbrace{\text{Tr}}_{\text{map from } L(B) \rightarrow \mathbb{C}}$.

$$\text{Tr}_B(\rho_A \otimes \rho_B) = \rho_A \cdot \text{Tr} \rho_B = \rho_A \text{ if } \rho_B \in \mathcal{S}(B).$$

Tr_B plays the role of taking the marginal distributions of a joint distribution.

$$\rho_{AB} = \sum_{a,b} P(a,b) |a\rangle\langle a| \otimes |b\rangle\langle b|$$

$\{|a\rangle\}_a$ ONB of A
 $\{|b\rangle\}_b$ ONB of B

then $\rho_A = \text{Tr}_B \rho_{AB} = \sum_a \left(\sum_b P(a,b) \right) |a\rangle\langle a|$

Check: • Complete ^{CP} positivity: $T \mapsto (I_A \otimes \langle b|) T (I_A \otimes |b\rangle)$ is CP

Sum of CP is CP.

• Trace-preserving: $\sum_b \text{Tr} (I \otimes \langle b| T I \otimes |b\rangle) = \sum_b \text{Tr} (I \otimes |b\rangle\langle b| T) = \text{Tr}(T).$

• Measurements.

$\{M_x\}_{x \in X}$ $\sum_x M_x^* M_x = I_A.$

Want the output of the channel to contain both the outcome x and the post measurement state.

Input: A

Output: $X \otimes A$ ^{post measurement state}.

\hookrightarrow holds outcome $\dim X = |X|$.
 fixed basis $\{|x\rangle\}_{x \in X}$

$$\mathcal{M} : L(A) \longrightarrow L(X \otimes A)$$

$$T \longmapsto \sum_{x \in X} \underbrace{|\alpha_x\rangle\langle\alpha_x|}_{\text{operator on } X} \otimes \underbrace{M_x T M_x^\dagger}_{\text{operator on } A}.$$

Check: • CP: $|\alpha_x\rangle\langle\alpha_x| \otimes M_x T M_x^\dagger = \underbrace{(|\alpha\rangle\langle\alpha| \otimes M_x)}_{L(A, X \otimes A)} T \underbrace{(\langle\alpha| \otimes M_x^\dagger)}_{\in L(X \otimes A, A)}$

\Rightarrow same argument.

• Trace-preserving:

$$\text{Tr}\left(\sum_x |\alpha_x\rangle\langle\alpha_x| \otimes M_x T M_x^\dagger\right) \stackrel{\substack{\downarrow \\ \text{Tr}(|\alpha_x\rangle\langle\alpha_x|) = 1}}{=} \sum_x \text{Tr}(M_x^\dagger M_x T) \stackrel{\substack{\downarrow \\ \sum_x M_x^\dagger M_x = I}}{=} \text{Tr}(T)$$

\uparrow
 $\text{Tr}(S \otimes T) = \text{Tr}(S) \cdot \text{Tr}(T).$

Can check that this models the measurement.

$$\mathcal{M}(\rho_A) = \sum_{x \in X} \underbrace{\text{Tr}(M_x \rho_A M_x^\dagger)}_{\text{prob of outcome } x} |\alpha_x\rangle\langle\alpha_x| \otimes \underbrace{\frac{M_x \rho_A M_x^\dagger}{\text{Tr}(M_x \rho_A M_x^\dagger)}}_{\text{post-measurement state conditioned on } x}.$$

Rk: Such a state is called a classical-quantum state.

i.e. of the form (cq state)

$$\rho_{XA} = \sum_{x \in X} P_X(x) |\alpha_x\rangle\langle\alpha_x| \otimes \rho_A^x$$

$\{|\alpha_x\rangle\}$ fixed basis, ρ_A^x density operator.

Representation of a quantum channel

3 ways of representing a quantum channel.

* Choi \leftarrow one operator in $L(A \otimes B)$.

* Kraus \leftarrow a list of operators in $L(A, B)$

* Stinespring \leftarrow an operator in $L(A, B \otimes E)$
 \leftarrow new space to be defined

The Choi operator: Fix a basis of A $\{|a\rangle\}_a$, let $\bar{A} \cong A$

$$J_{\bar{A}B}^{\mathcal{E}} = \sum_{a,a'} |axa'\rangle_A \otimes \mathcal{E}(|axa'\rangle_A) \in L(\bar{A} \otimes B)$$

$$= (\mathbb{I}_{\bar{A}} \otimes \mathcal{E})(\Phi)$$

Rk: We often just write A instead of \bar{A} . \leftarrow unnormalized maximally entangled state

Ex: • $\mathcal{E} = \mathbb{I}$ identity channel ($B \cong A$)

$$J_{\bar{A}B}^{\mathcal{E}} = \Phi = \sum_{a,a'} |axa'\rangle$$

• $\mathcal{E} = \text{Tr}$ ($B = \mathbb{C}$)

$$J_{\bar{A}B}^{\mathcal{E}} = \mathbb{I}_A$$

• $\mathcal{E}(S) = \text{Tr}(S)\sigma$ (constant output)

$$J_{\bar{A}B}^{\mathcal{E}} = \sum_{a,a'} |axa'\rangle \otimes \text{Tr}(\text{Tr}(|axa'\rangle) \sigma) = \mathbb{I}_A \otimes \sigma$$

Does $J^{\mathcal{E}}$ capture everything about \mathcal{E} ?

Choi-Jamiołkowski isomorphism : $L(L(A), L(B)) \rightarrow L(A \otimes B)$
 $\mathcal{E} \mapsto J^{\mathcal{E}}$

and its inverse is

$J \mapsto \left[S_A \mapsto \text{Tr}_A \left(\left(S_A^T \otimes I_B \right) J \right) \right]$
Transpose with respect to basis $\{|a_i\rangle\}$

Check:

$$\begin{aligned} \sum_{a, a'} |a\rangle\langle a'| \otimes \mathcal{F}(|a\rangle\langle a'|) &= \sum_{a, a'} |a\rangle\langle a'| \otimes \text{Tr}_A \left(\underbrace{(|a\rangle\langle a'|^T \otimes I)}_{|a'\rangle\langle a|} J \right) \\ &= \sum_{a, a'} |a\rangle\langle a'| \otimes \langle a'| (|a\rangle\langle a| \otimes I) J |a\rangle \\ &= \sum_{a, a'} |a\rangle\langle a'| \otimes \langle a| J |a\rangle \\ &= J. \end{aligned}$$

$J^{\mathcal{E}}$ can be used to easily check if \mathcal{E} is a valid quantum channel

Th: • $\mathcal{E} \in L(L(A), L(B))$ is completely positive

$J_{AB}^{\mathcal{E}} \geq 0$ \iff \mathcal{E} is trace-preserving $\iff J_A^{\mathcal{E}} = I_A$.
 $\text{Tr}_B \left(J_{AB}^{\mathcal{E}} \right)$

Conseq: • Complete positivity can be checked efficiently.

for all $(I_R \otimes \mathcal{E})(S) \geq 0$ for $S \in \text{Pos}(R \otimes A)$.

Theorem says sufficient to take $R \cong A$ and $S = \sum_{a, a'} |a\rangle\langle a'| \otimes |a\rangle\langle a|$

Proof: • \Downarrow is obvious

\Uparrow : $J_{AB}^{\mathcal{E}} = \sum_x \lambda_x |\psi_n \times \psi_n\rangle$ eigendecomposition.

Write $\mathcal{E}(S_A) = \sum_x K_x S_A K_x^*$, directly CP.

$$\begin{aligned} \mathcal{E}(S_A) &= \sum_x \lambda_x \text{Tr}_A \left((S_A^T \otimes I_B) |\psi_n \times \psi_n\rangle \right) \\ &= \sum_{x,a} \lambda_x \langle a | (S_A^T \otimes I_B) \cdot |\psi_n \times \psi_n\rangle |a\rangle \end{aligned}$$

Can write $|\psi_n\rangle = \sum_{a,b} \langle a,b | \psi_n \rangle |a\rangle_A \otimes |b\rangle_B$

So $|\psi_n \times \psi_n\rangle = \sum_{\substack{a,b \\ a',b'}} \langle a',b' | \psi_n \rangle \langle \psi_n | a,b \rangle |a' \times a\rangle \otimes |b' \times b\rangle$

$$\begin{aligned} \mathcal{E}(S_A) &= \sum_{\substack{x, a, a' \\ b, b'}} \lambda_x \langle a | S_A^T | a' \rangle |b' \times b\rangle \langle a', b' | \psi_n \rangle \langle \psi_n | a, b \rangle \\ &= \sum_{\substack{x \\ a, b, a', b'}} \sqrt{\lambda_x} \langle a', b' | \psi_n \rangle |b' \times b\rangle \langle a' | S_A | a \rangle \langle b | \langle \psi_n | a, b \rangle \cdot \sqrt{\lambda_x} \end{aligned}$$

$= \sum_x K_x S_A K_x^*$ where $K_x = \sum_{a',b'} \langle a', b' | \psi_n \rangle |b' \times a'\rangle \sqrt{\lambda_x}$

• \Rightarrow direct as $\text{Tr}(\mathcal{E}(|i \times j\rangle)) = \delta_{ij}$

$$\begin{aligned} \Leftarrow \text{Tr} \left(\text{Tr}_A \left((S_A^T \otimes I_B) J_{AB}^{\mathcal{E}} \right) \right) &= \text{Tr} \left((S_A^T \otimes I_B) J_{AB}^{\mathcal{E}} \right) \\ &= \text{Tr} \left(S_A^T \text{Tr}_B (J_{AB}^{\mathcal{E}}) \right) \\ &= \text{Tr} (S_A^T) \\ &= \text{Tr} (S_A) \quad \square \end{aligned}$$

Corollary: Any CP map $\mathcal{E}: L(A) \rightarrow L(B)$ can be written as

$$\mathcal{E}(S_A) = \sum_{\alpha=1}^r K_{\alpha} S_A K_{\alpha}^*$$

where $K_{\alpha} \in L(A, B)$ \leftarrow called Kraus operators

with $r \leq (\dim A)(\dim B)$ \leftarrow actually $r = \text{rank}(J^{\mathcal{E}})$.

\mathcal{E} is trace preserving iff $\sum_{\alpha} K_{\alpha}^* K_{\alpha} = I_A$

Rk: Called **operator-sum representation**,
Kraus

(Stinespring dilation)

Corollary: Any CP map $\mathcal{E}: L(A) \rightarrow L(B)$

can be written as

$$\mathcal{E}(S_A) = \text{Tr}_E(M S_A M^*)$$

where $M \in L(A, B \otimes E)$

with $\dim E \leq (\dim A)(\dim B)$

\mathcal{E} is trace-preserving iff

$$\underline{M^* M = I_A}$$

M is an isometry

ie. preserves norms

$$\|M|\psi\rangle\| = \|\psi\rangle\|$$

Proof:

$$\text{write } \mathcal{E}(S_A) = \sum_{\alpha} K_{\alpha} S_A K_{\alpha}^*$$

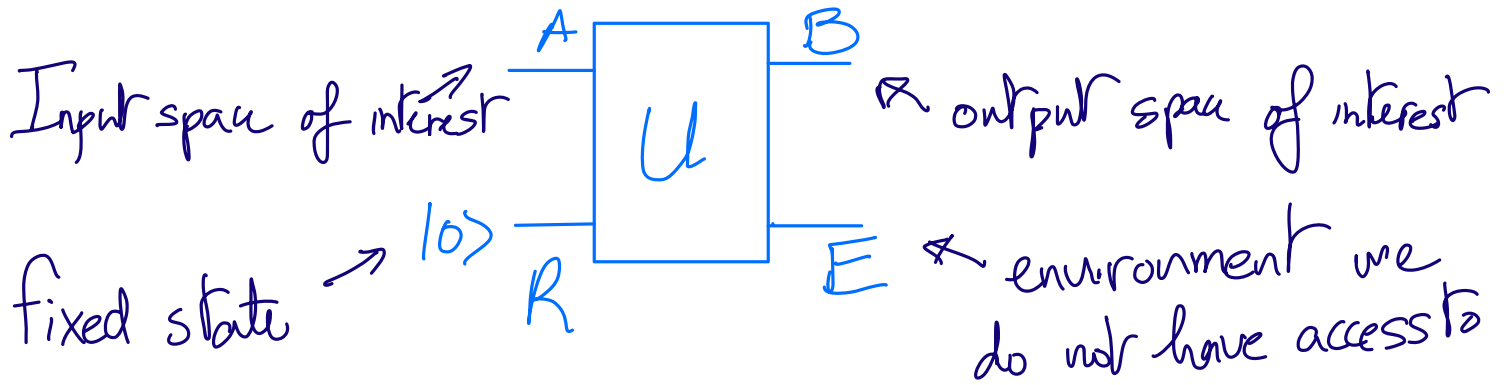
let $E = \text{span}\{|\alpha\rangle\}$

and $M = \sum_i K_i \otimes |i\rangle$

Then $MS_A M^\dagger = \sum_{x, x'} K_x S_A K_{x'}^\dagger \otimes |x\rangle\langle x'|_E$ \square

Interpretation:

Can see any evolution modeled by quantum channel $\mathcal{E}: L(A) \rightarrow L(B)$ as a **unitary** evolution



$U: |\psi\rangle_A \otimes |0\rangle_R \mapsto M|\psi\rangle$ (check: such a unitary exists)

A concept we will use: operator monotone convex functions

$f: I \rightarrow \mathbb{R}$ operator monotone if for any $d \geq 1$

for any Hermitian operators $A, B \in \text{Herm}(\mathbb{C}^d)$
with $\text{spec}(A), \text{spec}(B) \subseteq I$.

$$\underbrace{A \geq B}_{\text{means } A-B \in \text{Pos}(\mathbb{C}^d)} \Rightarrow f(A) \geq f(B)$$

Rk: There exist functions that are monotone but not operator monotone e.g. $x \mapsto x^2$ (Ex: find example)

But $x \mapsto \log x$ & $x \mapsto \sqrt{x}$ are operator monotone.

$f: I \rightarrow \mathbb{R}$ operator convex if for any $d \geq 1$

for any Hermitian operators $A, B \in \text{Herm}(\mathbb{C}^d)$
with $\text{spec}(A), \text{spec}(B) \subseteq I$.

$$f(\lambda A + (1-\lambda)B) \leq \lambda f(A) + (1-\lambda)f(B).$$

Rk: . Example of non-operator convex: $x \mapsto x^3$.

. Example operator convex: $-\log x, \frac{1}{x}$